



**ERICSSON**

# OPTIMIZING TCP FOR RADIO

# RAN CHARACTERISTICS



## › Congestion

- Increasing number of users and/or increasing traffic in a cell
- Lower possibility of overprovisioning than in fixed networks

## › Reduced signal strength

- when the device moves towards the edge of the cell, the Signal to Interference plus Noise Ratio (SINR) decreases

## › Handovers

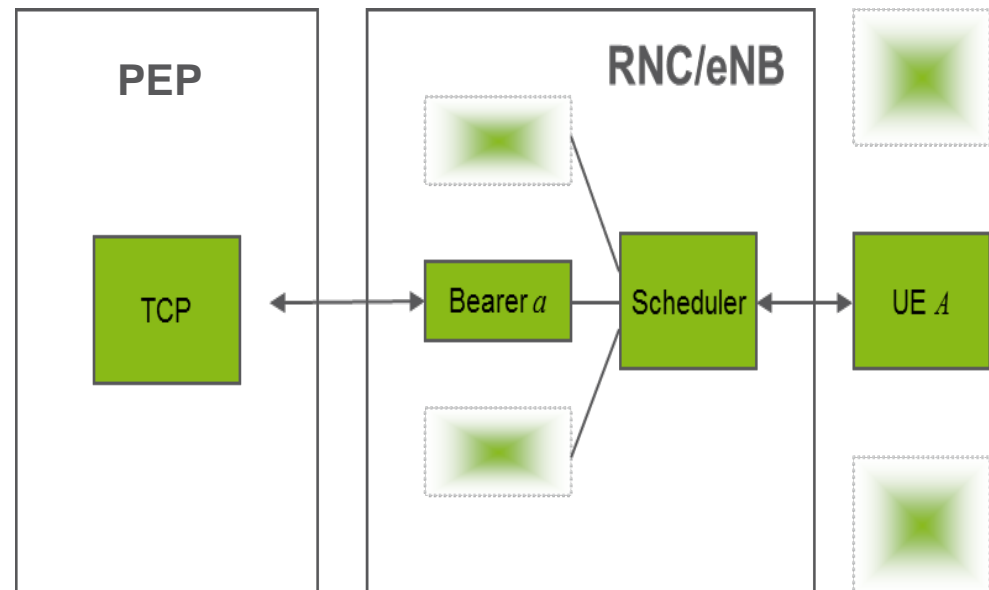
- Between cells
- Between radio technologies

## › Varying channel conditions mean that it makes sense to wait.

# RAN TRANSPORT REQUIREMENTS



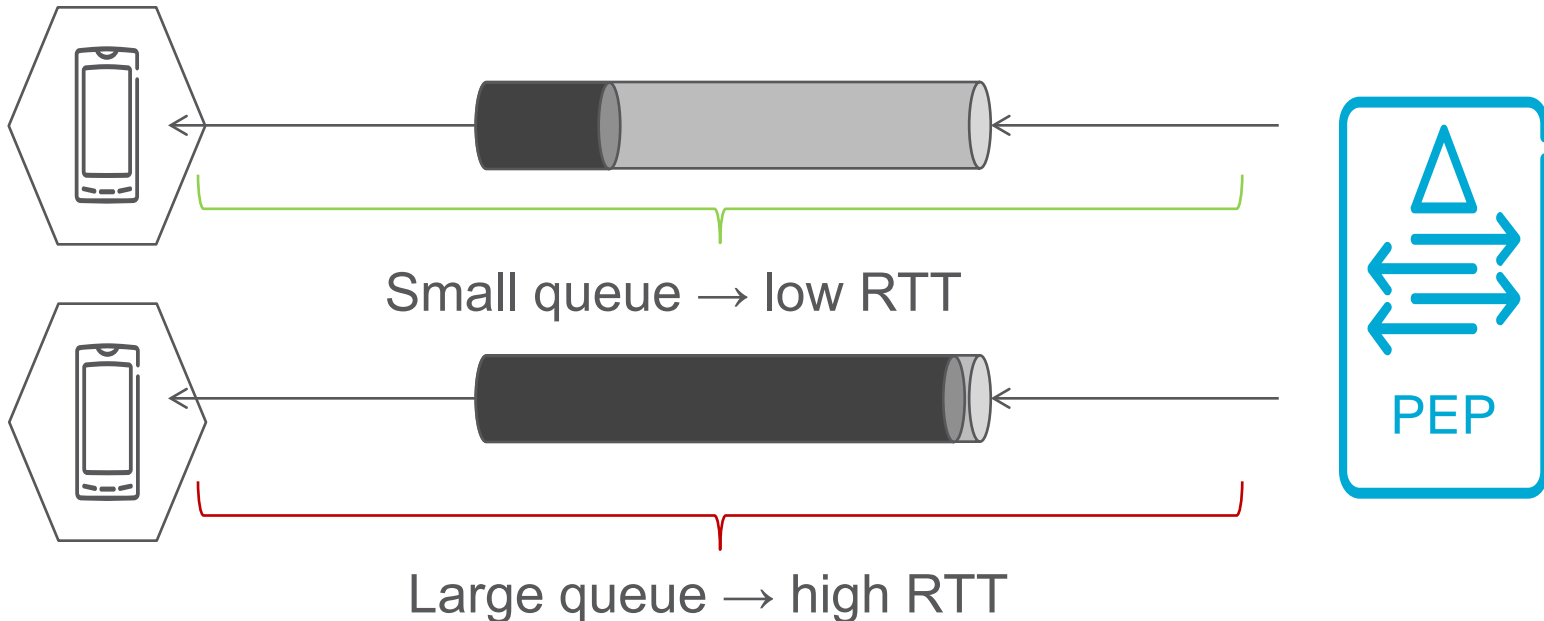
- › Keep the RNC/eNB scheduler queues non-empty.
  - Unnecessarily empty buffers may reduce throughput.
- › Keep the scheduler queues short.
  - Bloating buffers decrease responsiveness, increase waste when flows are abandoned and add delay to handovers.
- › Balance flow rates to optimise user experience.
  - Higher throughputs for interactive flows than for background flows.



# THEORY AND MODELS



- $\uparrow$ Users in a cell  $\rightarrow$   $\downarrow$ Individual capacity per user  $\rightarrow$   $\uparrow$ Queues in RNC/eNB
- $\uparrow$ Distance from a base station  $\rightarrow$   $\downarrow$ Signal Strength  $\rightarrow$   $\uparrow$ Queues in RNC/eNB
- $\uparrow$ Queue length  $\rightarrow$   $\uparrow$  RTT per individual flow



# TCP TRANSMISSION RATE



- › TCP is designed to "fill the pipe".
  - Filled pipe → packet drops → TCP reduces transmission rate
- › TCP losses are "feared" by router designers
  - Modern equipment often comes with large buffers – including RNC/eNB
- › TCP fills those buffers with packets.
  - Buffer bloat problematic for new or interrupted flows and handovers.

# TCP OPTIMIZATION



- › TCP optimization helps to mitigate congestion by protecting the network from aggressive traffic.
- › During low to medium cell load, it is in general beneficial to push data fast.
- › But when there are limitations in the radio network, aggressive TCP behavior can fill the buffers, provoke large loss rates and make TCP slow down considerably.
- › End user pays for retransmitted data.

# RADIO OPTIMIZED HYBRID CCA



- › Treats both packet loss and increased RTT as congestion signal
- › Separate back-off factors for loss and increased latency
- › Sensitivity to RTT increase and back-off factors configurable per flow
- › Configurable ICW

# SINGLE POINT OF CONTROL



- › Hybrid or delay sensitive algorithms typically yield to pure loss based flows.
- › A single point of in the operator network that performs congestion control can ensure that overall fairness between flows.
- › Initial window size can be auto-tuned based on knowledge about current conditions for specific user. TCP needs less time to probe the network.
  
- › Simple to deploy

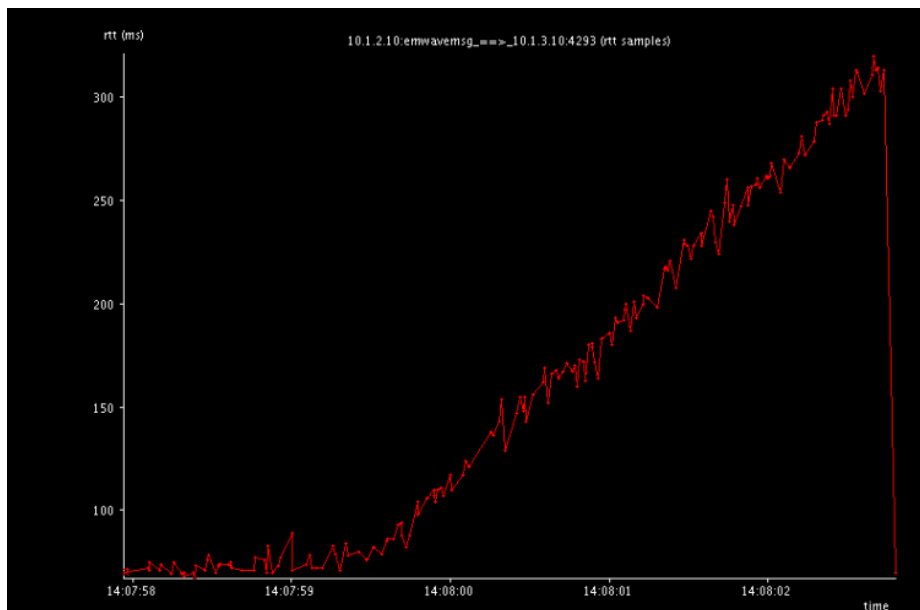


# CCA EFFECTS ON RTT

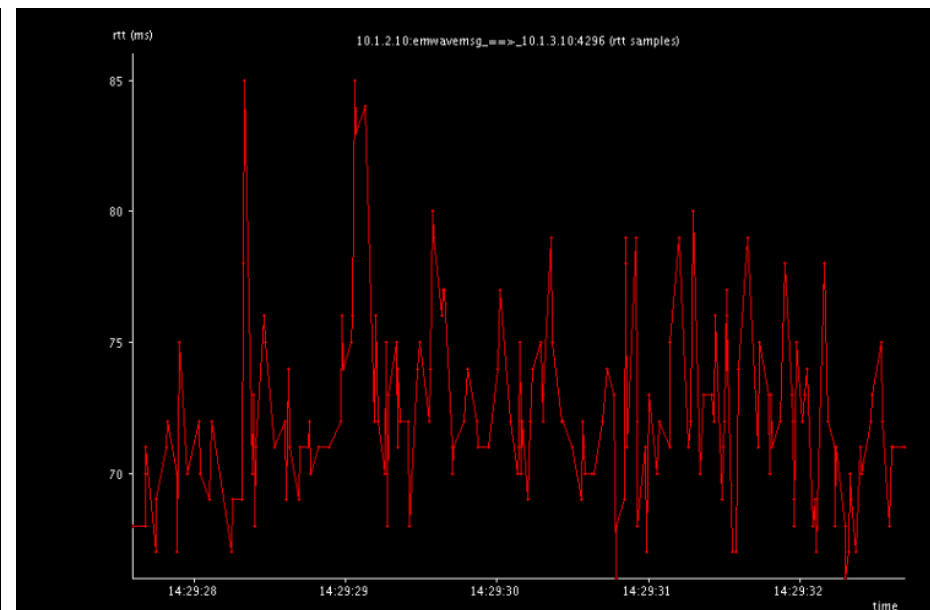


Artificial example without proper AQM

New Reno



Ericsson Hybrid CCA

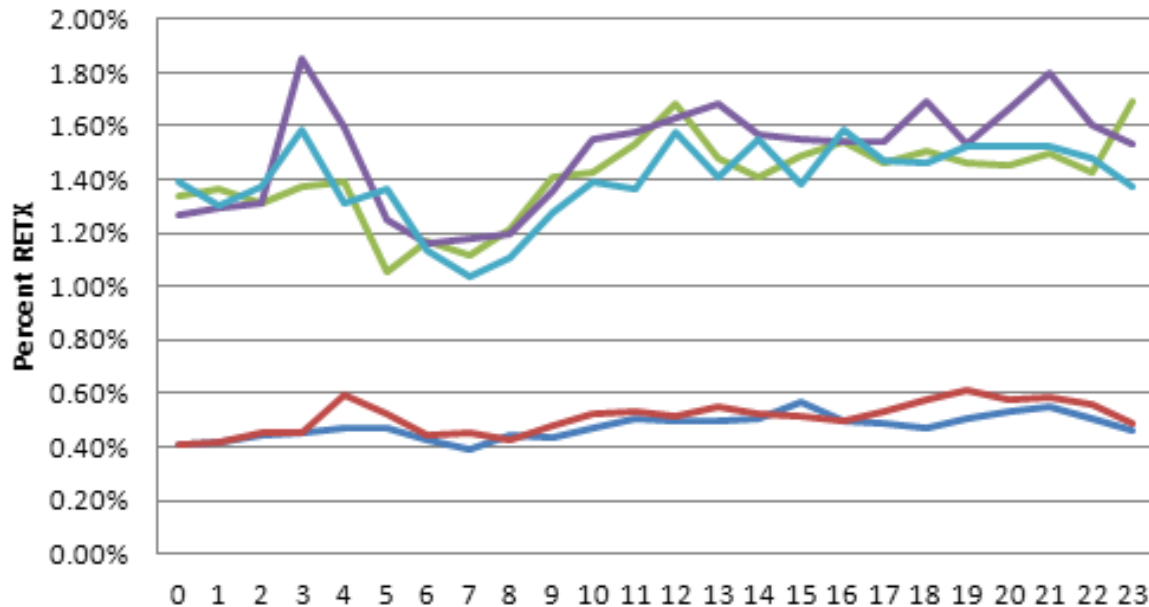


Note different scales!

# LIVE DATA - RETRANSMISSIONS



Data collected at large site of Tier 1 operator in emerging market.  
Up to 70% reduction of retransmissions after introducing new  
CCA.



# DYNAMIC ICW – RESULTS



Well dimensioned LTE network, European Operator.

Measurements on page load times for most popular sites in the country

Page load times reduced by 5 - 15% by using dynamic ICW.

# RELATION TO ENCRYPTION & WAY FORWARD



- › Current best practice is to deploy a “transparent” proxy outside the packet gateway.
- › Transport protocols are being encrypted basically bypassing the PEP in its current form.
- › Can we make a PEP explicit, how would we go on doing that?
  - Reverse proxy, origin consent?
  - Forward proxy, client consent?
- › A PEP is deployed within an operator domain, thus simplifying the trust relationship, potentially allowing richer signaling.



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